

lecture 1: see slides
 lectures 2-5: handwritten notes

Real applied multigraded commutative algebra

24th National School on Algebra
 EMS Summer School on Multigraded Algebra & Applications
 Mistral Resort, Moieciu de Sus, Romania
 17-24 August 2016, Thursday 11:10-12:00
 Friday 9:00-9:50
 Saturday 11:10-12:00, 12:10-13:00
 Monday 10:00-10:50

$x = x_1, \dots, x_n$

k field

$k[x] = \bigoplus_{b \in \mathbb{N}^n} k \cdot x^b$

\mathbb{Z}^n -graded

module M \mathbb{Z}^n -graded if $M = \bigoplus_{a \in \mathbb{Z}^n} M_a \iff M$ is \mathbb{Z}^n -module/ k .

$x^b: M_a \rightarrow M_{a+b}$

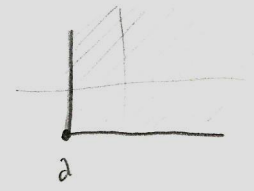
M free/ $k[x]$

rank 1: basis $e \in M_a$



$\Rightarrow M \hookrightarrow k[\mathbb{Z}^n]$

$k[x] \cdot x^a = k[x](-a)$



rank ≥ 1 : \bigoplus rank 1.

now you're probably used to seeing it

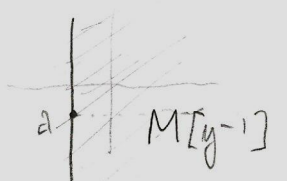
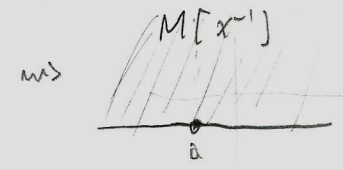
localization

$M[f^{-1}]$ \mathbb{Z}^n -graded needs f homogeneous
 $\iff f = x^b$ some $b \in \mathbb{N}^n$

may as well $b \in \{0,1\}^n \iff \sigma \subseteq \{1, \dots, n\} = [n]$.

M has only 2^n \mathbb{Z}^n -graded localizations $M[x^{-\sigma}]$

e.g. $M = k[\mathbb{Z}^n_{\geq a}] \Rightarrow M[x^{-\sigma}] = k[a + \mathbb{Z}^{\sigma} \times \mathbb{N}^{\bar{\sigma}}]$
 $= k[a + \mathbb{N}^n]$

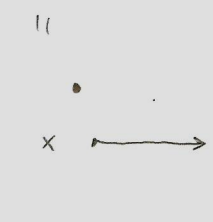
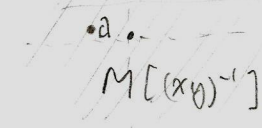


[local duality, 3.5]

Lemma: $M \cong \bigoplus$ localizations of free mods

$\Downarrow \Uparrow$ if M is \mathbb{Z}^n -finite: $\dim_k M_a < \infty \forall a \in \mathbb{Z}^n$

M flat
 (in \mathbb{Z}^n -mod \Rightarrow in $k[x]$ -mod)



Matlis duality

intuitively: turn M upside-down

precisely: Def: The Matlis dual of M is M^\vee with

$$(M^\vee)_a = (M_{-a})^*$$

$(-)^* = \text{Hom}_k(-, k)$ vector space dual.

Prop^v: $M \mapsto M^\vee$ is exact.

$M \xrightarrow{\sim} (M^\vee)^\vee$ if M \mathbb{Z}^n -finite.

M flat $\Leftrightarrow M^\vee$ injective

$(M \otimes - \text{ exact}) \Leftrightarrow (\text{Hom}(-, M) \text{ exact})$

Pf: Exercise!

e.g. $k[x][x^{-\sigma}]^\vee = \left(k[\mathbb{Z}^\sigma \times \mathbb{N}^{\bar{\sigma}}]^\vee \right) = R$

$k[\mathbb{Z}^\sigma \times \mathbb{N}^{\bar{\sigma}}] = E(k[x]/\mathfrak{A}_\sigma) = E(k[\mathbb{N}^\sigma])$
 $\mathfrak{A}_\sigma = \langle x_i \mid i \notin \sigma \rangle$

Def: indecomposable injective



Prop [GW '78] $M \cong \bigoplus$ indec. injectives

$M \xleftrightarrow{\sim} M^\vee$ injective (in \mathbb{Z}^n -mod) (no \mathbb{Z}^n -finiteness required.)

Pf: Same as for R -mod, R noeth.

Cor: Lemma

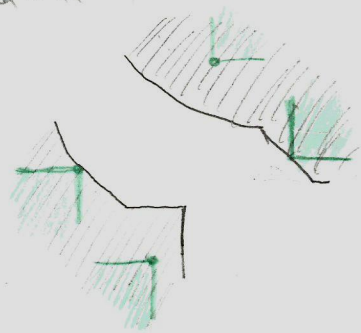
Pf: Prop^v: $M \text{ flat} \Rightarrow M^\vee \text{ inj} \Rightarrow M^\vee \cong \bigoplus \Rightarrow M^{\vee\vee} = \prod M$
 $\stackrel{||}{=} M$ if \mathbb{Z}^n -finite. \square

Analogues for arbitrary posets \mathcal{Q} ?

e.g. \mathcal{Q} finite (\Rightarrow flat is hard to do)

$\mathcal{Q} = \mathbb{R}^n \Rightarrow$

Def: $U \subseteq \mathcal{Q}$ upset if $Q \succeq u \in U \forall u \in U$
 $D \subseteq \mathcal{Q}$ downset if $Q \preceq d \in D \forall d \in D$.



We'll be decomposing \mathcal{Q} -mods using these instead of tree, flat, or inj.

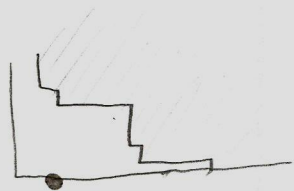
$S \subseteq \mathcal{Q} \Rightarrow k[S] = \bigoplus_{s \in S} k_s$

Lemma: $k[U] \subseteq k[\mathcal{Q}]$ indicator submodule
 $k[\mathcal{Q}] \twoheadrightarrow k[D]$ indicator quotient module

e.g. indec. free or flat \Rightarrow indicator submod.

" injective \Rightarrow " quotient

monomial ideal in $k[x]$ = indicator submod of upset in \mathbb{N}^n



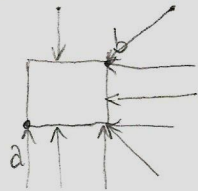
To lift from \mathbb{Z}^n to \mathbb{R}^n or \mathcal{Q} , need:
 or \mathcal{Q}

Def: $\pi: \mathcal{Q} \rightarrow \mathcal{P}$ is a poset morphism if $q \preceq q' \Rightarrow \pi(q) \preceq \pi(q')$.

cf. Jürgen's lecture yesterday

e.g. \mathcal{Q} finite $\Leftrightarrow \mathbb{Z}^n = \mathcal{P}$ n minimal $\Rightarrow n = \text{odim } \mathcal{Q}$

$\mathcal{Q} = \mathbb{Z}^n \twoheadrightarrow \mathcal{P} = [a, b]$ a box
 \uparrow
 convex projection



• upset $U = \pi^{-1}(\{1 \in \{0,1\}\})$

• downset $D = \pi^{-1}(\{0 \in \{0,1\}\})$

Lemma: $\pi^{-1}(\text{upset}) = \text{upset}$

$\pi^{-1}(\text{downset}) = \text{downset}$

cf. Jürgen's lecture yesterday

Ex: $\text{Hom}_{\mathcal{Q}\text{-mod}}(k[U], k[U']) = ?$

$k[D], k[D]$

$k[U], k[D]$

$k[D], k[U]$

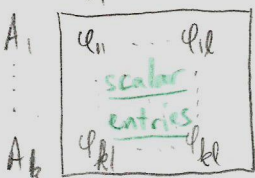
DO ON 5

Monomial matrices

column labels

objects

B_1, \dots, B_l



monomial matrix

representing if $\varphi_{ij} = 0$ whenever Hom dictates it has to

row labels

$A_1 \oplus \dots \oplus A_k$

$\longrightarrow B_1 \oplus \dots \oplus B_l$

if $\text{Hom}(A_i, B_j) = k$ or 0

A

\longleftarrow

B

B_j, A_i

e.g. $A \rightarrow B$ both \mathbb{Z}^n -finite injective

$A_i = k[D_i], B_j = k[D_j]$

$A \leftarrow B$ " " flat

U_i, U_j

$A \rightarrow B$ or $A_i = k[U_i], B_j = k[D_j]$

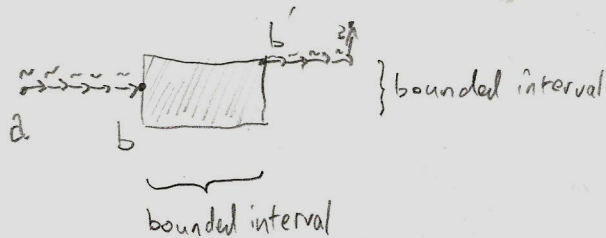
flat injective

Def: \mathbb{Z}^n -finite M is finitely determined if $\forall i=1, \dots, n$

$$\cdot x_i: M_a \xrightarrow{\sim} M_{a+e_i}$$

whenever $a_i \notin$ some bounded interval.

"Everything about M is determined by what happens inside of a box"



e.g. M f.g. = all gens $\geq b$

all relations $\leq b'$

e.g. \oplus localizations of $k[x](-a) = \mathbb{Z}^n$ -finite flat

e.g. M f.d. $\Leftrightarrow M^v$ f.d.

e.g. \mathbb{Z}^n -finite injective

Def: Fix \mathbb{Z}^n -mod M .

1. injective hull of M is $M \hookrightarrow E$ with E injective
(envelope)

- finite if E \mathbb{Z}^n -finite
- minimal if # summands of E is minimal.

2. injective resolution of M is a complex

$$E^\bullet: 0 \rightarrow E^0 \rightarrow E^1 \rightarrow \dots \rightarrow E^l \rightarrow \dots$$

of injectives E^i with $H^i E^\bullet = \begin{cases} M & i=0 \\ 0 & i \geq 1 \end{cases}$

- length l if $E^i = 0$ for $i > l$ and $E^l \neq 0$.
- finite if E^\bullet is \mathbb{Z}^n -finite
- minimal if # summands of E^i is minimized $\forall i$.

$\Leftrightarrow M \rightarrow E^0$ and $\text{coker}(E^{i-1} \rightarrow E^i) \rightarrow E^{i+1}$ are injective hulls $\forall i \geq 1$.

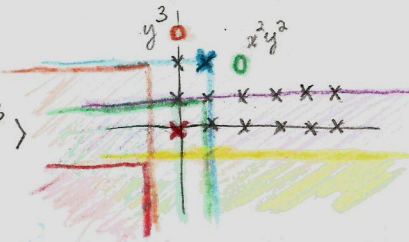
Thm ^{f.d.}: M f.d. $\Leftrightarrow M$ has ^{finite} minimal inj. res. of length $\leq n$,
unique up to \cong .

Pf: Basically [GW '78]. Essential point: $\text{ann}(y) = \emptyset \Rightarrow M \xrightarrow{y \in M_2} \mathbb{k}[Z + \mathbb{Z}^\sigma - \mathbb{N}^\sigma]$
with $\psi(y) \neq 0$. Primary decomp. (of principal submods) \Rightarrow every elt. of M has multiple
that is one of finitely many such cogenerators. \square

Cor: Same for flat resolutions! [$z \in M_2 \Rightarrow z' \in M_2'$ with " $z \cong z'$ " and $\alpha \in [b, b']$; $M_{[b, b']}$ f.g.]

Pf: (Thm)^v.

E.g. $M = \mathbb{k}[x, y] / \langle x^2 y, y^3 \rangle$



$$0 \rightarrow \mathbb{k}[0, 2] - \mathbb{N}^2 \oplus \mathbb{k}[0, 1] - \mathbb{Z} \times \mathbb{N} \rightarrow \mathbb{k}[1, 0] - \mathbb{N}^2 \oplus \mathbb{k}[0, -1] - \mathbb{Z} \times \mathbb{N} \oplus \mathbb{k}[-1, 2] - \mathbb{N}^2 \rightarrow \mathbb{k}[-1, 1] - \mathbb{N}^2 \rightarrow \dots$$

Monomial matrices

$$\begin{bmatrix} 1 & 0 & 1 \\ 1 & 1 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 1 \\ -1 \\ -1 \end{bmatrix}$$

from p. (4)

$$+ \text{Hom}_{\mathbb{Q}\text{-mod}}(-, -) = \mathbb{k}$$

from p. (3)

Exercise: Matlis dual of this monomial matrix.

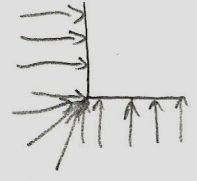
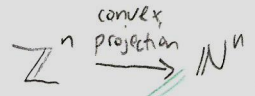
Poset encoding

Def: Fix $\pi: Q \rightarrow P$ and P -mod M . The pullback of M along π is

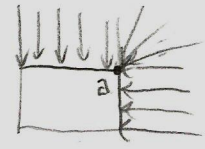
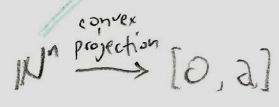
Q -mod $\pi^* M$, which has $(\pi^* M)_q = M_{\pi(q)}$.

e.g. $k[U] = \pi^* \begin{pmatrix} k \\ 0 \end{pmatrix}$

$k[D] = \pi^* \begin{pmatrix} 0 \\ k \end{pmatrix}$



$\Rightarrow \pi^* M = \check{C}M$ Čech hull



$\Rightarrow \pi^* M = P_a M$ positive extension

e.g. M is squarefree if $M = P_{(0,1)} M$

M f.d. $\Leftrightarrow M = \pi^* M_{[a,b]}$ for some $a \leq b$

Important: insert Functoriality, p. 8

$\pi: \mathbb{Z}^n \rightarrow [a,b]$ convex projection

Def: Fix poset Q and Q -mod H .

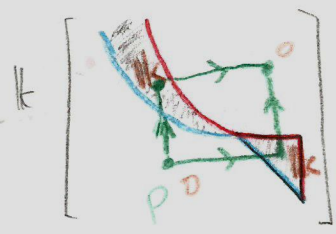
An encoding of H by a poset P is

- morphism $\pi: Q \rightarrow P$ and
- P -mod M with $H \cong \pi^* M$.

The encoding is finite if

- $|P| < \infty$ and
- $\dim_k M_p < \infty \forall p \in P$.

e.g. all the examples we just saw, but even more powerful as (finite) data structure for modules/ \mathbb{R}^n .



$\cong \pi^* H$
 $\pi: \mathbb{R}^2 \rightarrow P$

Indicator resolutions

Def: Fix \mathcal{Q} -mod. H .

1. F_\bullet a homological indicator resolution of H if

- $F_i = \bigoplus (\text{indicator submods of } k[\mathcal{Q}]) \quad \forall i$

- $\dots \leftarrow F_{i-1} \leftarrow F_i \leftarrow \dots$

- $H_i F_\bullet = \begin{cases} 0 & i \neq 0 \\ H & i = 0 \end{cases}$

$$F_\bullet: \dots \leftarrow F_{i-1} \leftarrow F_i \leftarrow \dots$$

2. E^\bullet a cohomological indicator resolution of H if

- $E^i = \bigoplus (\text{indicator quotients of } k[\mathcal{Q}]) \quad \forall i$

- $\dots \rightarrow E^i \rightarrow E^{i+1} \rightarrow \dots$

- $H^i E^\bullet = \begin{cases} 0 & i \neq 0 \\ H & i = 0 \end{cases}$

$$E_\bullet: \dots \rightarrow E^i \rightarrow E^{i+1} \rightarrow \dots$$

3. length: $\max \{i \mid F_i \neq 0\}$

$\max \{j \mid E^j \neq 0\}$

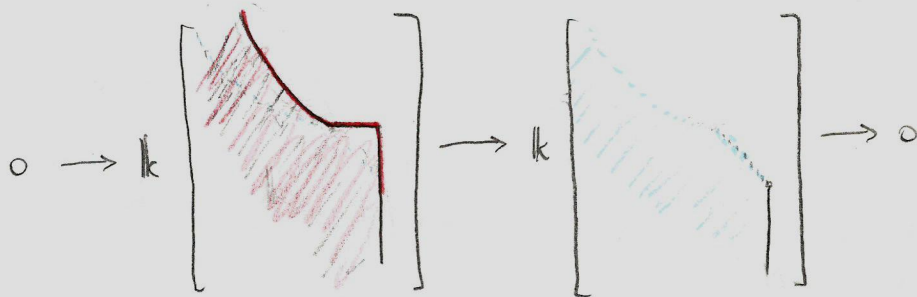
$\Rightarrow \text{length} < \infty$


4. finite if #indicator summands $< \infty$

5. subordinate to an encoding of H if each summand is constant on every fiber of the encoding

\Leftrightarrow Each summand is pulled back from an indicator P -module.

e.g.

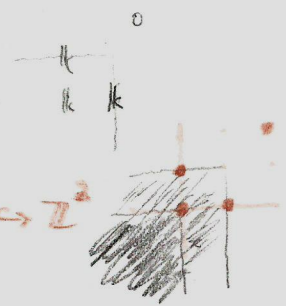


subordinate to encoding by  from before. (Why? fibers of $\pi \dots$)

e.g. Thom f.d.!

at "see previous e.g."

$$= \pi^* (0 \rightarrow k \left[\begin{array}{c} \cdot \\ \cdot \\ \cdot \end{array} \right] \rightarrow k \left[\cdot \right] \rightarrow 0)$$



$$= \pi^* \pi^* (0 \rightarrow k \left[\begin{array}{c} \cdot \\ \cdot \\ \cdot \end{array} \right]^{[1]} \rightarrow k \left[\begin{array}{c} \cdot \\ \cdot \\ \cdot \end{array} \right]^{[0]} \rightarrow 0) \text{ for } P \hookrightarrow \mathbb{Z}^2$$

Syzygy Thm: Fix poset \mathcal{Q} and \mathcal{Q} -mod H with ^{finite} encoding $\pi: \mathcal{Q} \rightarrow P$.
 $H \cong \pi^* M$.

1. H admits finite homological indicator res. subordinate to π .
2. " " " "

Prop: Suppose $i: P \subseteq Z$. $\exists Z$ -mod $i_* M$ (pushforward) with
 poset morphism
 think $Z = \mathbb{Z}^m$

• $i_* M|_P = M$ and is universal with this property:

• $i_* M \rightarrow N$ whenever $N|_P = M$.

Pf: At $z \in Z$, $(i_* M)_z = \varinjlim_{p \leq z} M_p$. "Kan extension" \square

Pf of Thm: $|P| < \infty \Rightarrow i: P \hookrightarrow \mathbb{Z}^m$ some m (order dim of P).

$N = i_* M$ f.d. because $|P| < \infty$: only finitely many \varinjlim 's, and they only change when some coord. of z passes that coord. of some $i(p)$.

N has finite flat or inj. res. by Thm f.d.

$i^*(\text{flat res } N) = \text{homological indicator res of } M$

inj. = co

$\pi^*(\text{res of } M) = \text{res of } H \text{ subordinate to } \pi$. \square

e.g. see previous e.g.

Functoriality (see p. 6)

Prop: π^* is functorial: $M \rightarrow M'$
 $\Rightarrow \pi^* M \rightarrow \pi^* M'$.

think A_i 's and B_j 's are downsets

On monomial matrices, if $M = k[A_1] \oplus \dots \oplus k[A_k] \rightarrow k[B_1] \oplus \dots \oplus k[B_l] = M'$
 or $M' \leftarrow M$ if A_i 's and B_j 's are upsets

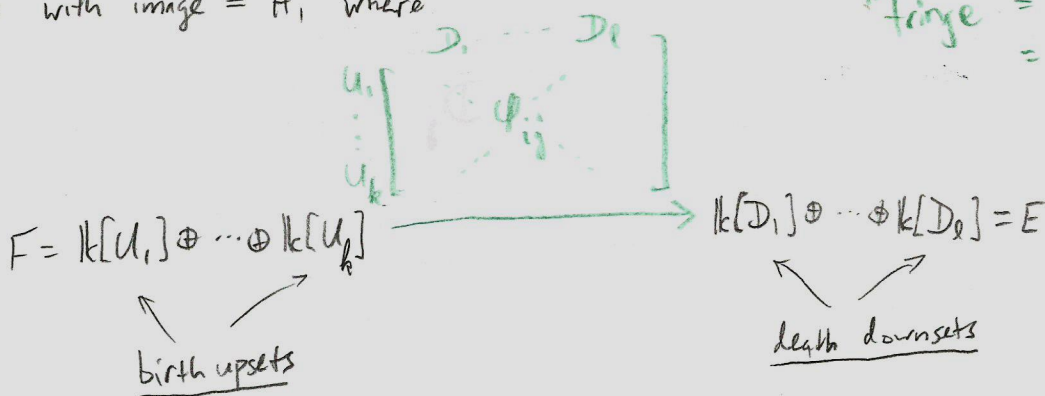
$$\pi^* \left(\begin{matrix} B_1 & \dots & B_l \\ \left[\begin{matrix} A_1 \\ \vdots \\ A_k \end{matrix} \right] \cdot \left[\begin{matrix} \phi_{ij} \end{matrix} \right] \end{matrix} \right) = \begin{matrix} \pi^{-1}(B_1) & \dots & \pi^{-1}(B_l) \\ \left[\begin{matrix} \pi^{-1}(A_1) \\ \vdots \\ \pi^{-1}(A_k) \end{matrix} \right] \cdot \left[\begin{matrix} \phi_{ij} \end{matrix} \right] \end{matrix}$$

or ... \uparrow

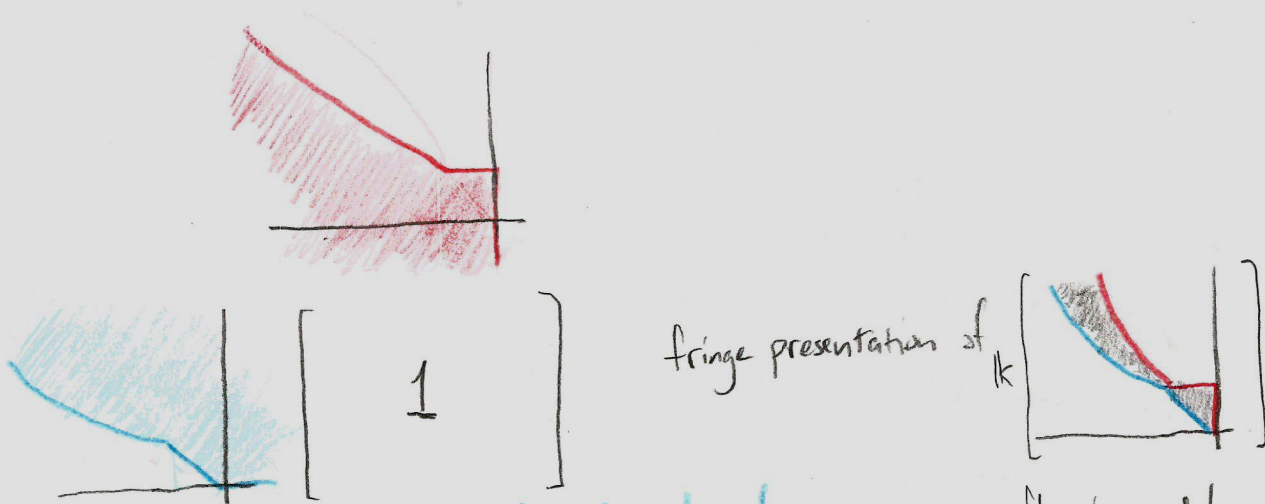
Fringe presentation

Def: A fringe presentation of a \mathbb{Q} -mod. H is a homomorphism

$F \rightarrow E$ with $\text{image} \cong H$, where

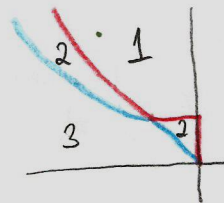


e.g.



topological interpretation: classes begin showing up at lower boundary of blue parameters and cease to appear at upper boundary of red parameters.

It's wing module, essentially, though there it's really Hilbert function



Cor: Every finitely encoded \mathbb{Q} -module admits fringe presentation.

Pf: Like O_n map in Tate resolution:

$$\text{coker}(F_1 \rightarrow F_0) = \ker(E^0 \rightarrow E^1). \square$$

Importance: interpretable data structure for statistics on sets of persistence modules

Problem for stats [Carlsson-Zomorodian 2008] (though lots of us here aren't at all surprised by this)

Even f.g. \mathbb{Z}^n -mods have moduli of $\dim \geq 1$: scalar entries ϕ_{ij} can vary continuously $\circ!$

So how do you stats on sets of these?

Answer 1: geometric probability and stats on samples from stratified spaces - a completely different lecture series

- LLN
- CLT
- PCA ...

Answer 2: "no-moduli" conj. for fly wings persistence in dim 0 } this is all there is!
 codim 1

⇒ fixed $[e_{ij}]$; only pay attention to

- moving U_i and D_i .
- ⇔ • rank function $(Q \times Q)_{\pm \Delta} \rightarrow \mathbb{N}$
- $\{(g, g') \mid g \neq g'\}$

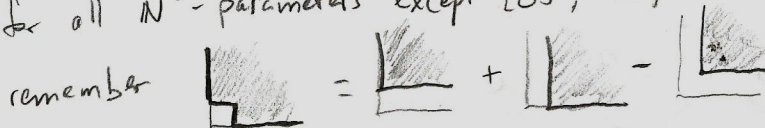
fringe presentation and data structure for (semianalytic) functions on $(\mathbb{R}^n \times \mathbb{R}^m)_{\pm \Delta}$ of this sort: level sets are semialgebraic (for fly wings and most applications)

- compute
- store
- manipulate/analyze statistically


Note: even for f.g. \mathbb{Z}^n -mods, fringe presentation is


- new
- more efficient than free presentation:

topologically speaking, if a class appears for all \mathbb{N}^2 -parameters except $[0]$, why



Instead, ignore "obvious" relations/syzygies that collapse/gather arise from upsets/downsets not being principal.

Exercise: $X =$ 

x-filtration: 

y-filtration: 

compute: gens + rels

- fringe presentation

for filtration of X by \mathbb{N}^2 .

Existence

You can and should ask me: your thms all assume existence of finite encoding;
When does that happen? Is existence restrictive?

Def: Fix \mathcal{Q} -mat. H .

H induces isotypic regions: classes for equivalence relation generated by $a \sim b$ whenever

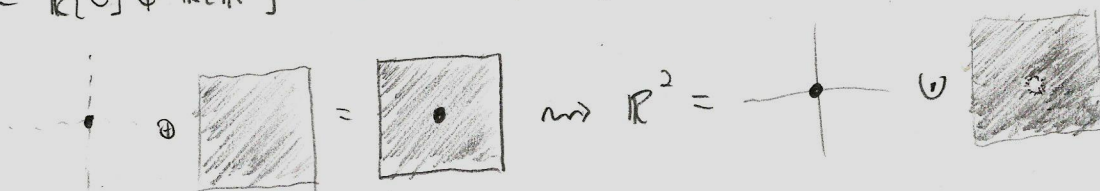
- $\bullet a \leq b$ in \mathcal{Q}
- $\bullet H_a \xrightarrow{\sim} H_b$.

H is tame if

- $\bullet \dim_{\mathbb{k}} H_g < \infty \quad \forall g \in \mathcal{Q}$
- $\bullet \# \text{ isotypic regions} < \infty$

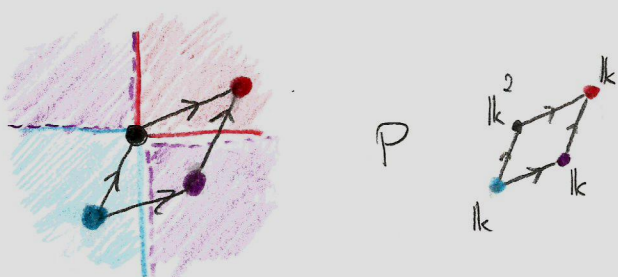
} always happens in data analysis

e.g. $H = \mathbb{k}[0] \oplus \mathbb{k}[\mathbb{R}^2]$ induces isotypic regions $\{0\}$ and $\mathbb{R}^2 \setminus \{0\}$.

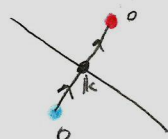


Encoding Thm: tame $\Rightarrow \exists$ finite encoding.

e.g. No poset map $\mathbb{R}^2 \rightarrow \mathcal{P}$ has fibers $\{0\}$ and $\mathbb{R}^2 \setminus \{0\}$.

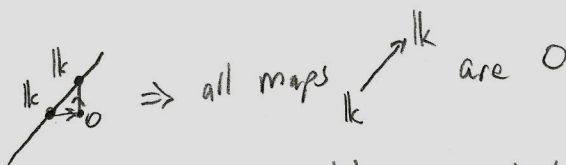


e.g. 1. $H = \mathbb{k} \left[\begin{array}{c} \diagdown \\ \diagup \end{array} \right]$ over \mathbb{R}^2 :



$a < b \Rightarrow H_a \rightarrow H_b$ is 0.
 \Rightarrow uncountably many isotypic regions, but: finite encoding.

2. $H = \mathbb{k} \left[\begin{array}{c} \diagup \\ \diagdown \end{array} \right]$ over \mathbb{R}^2 :



\Rightarrow all maps $\mathbb{k} \rightarrow \mathbb{k}$ are 0
 \Rightarrow uncountably many isotypic regions and no finite encoding

Conclusion: we think "finitely encoded" is the right notion. \square